

CSIRO DIVISION OF RADIOPHYSICS

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AT FILE NOTE

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RNM

POSSIBLE STATION LOCATIONS FOR COMPACT ARRAY

The table below gives station positions computed by Geoff Poulton for the 6 km compact array in a configuration based on 20 metre increments for 1.5, 3 and 6 km baselines. This layout is interim in a sense as we are looking at further optimization of the configuration perhaps involving non-grating arrays. It does, however, represent a practical solution in that future solutions are likely to have a similar number of and overall distribution of stations - in particular the main concentration at the western end of the 3 km track and two stations at the 6 km point.

<u>Stn.</u>	<u>Intervals</u> (20m)	<u>Distance</u> (m)	<u>Stn.</u>	<u>Intervals</u> (20m)	<u>Distance</u> (m)	<u>Stn.</u>	<u>Intervals</u> (20m)	<u>Distance</u> (m)
1	0	0	13	77	1540	25	124	2480
2	2	40	14	79	1580	26	125	2500
3	4	80	15	81	1620	27	134	2680
4	8	160	16	83	1660	28	135	2700
5	20	400	17	88	1760	29	138	2760
6	24	480	18	94	1880	30	140	2800
7	30	600	19	96	1920	31	142	2840
8	41	820	20	97	1940	32	146	2920
9	50	1000	21	102	2040	33	148	2960
10	57	1140	22	106	2120	34	149	2980
11	64	1280	23	112	2240	35	150	3000
12	76	1520	24	118	2360	36	296	5920
						37	300	6000

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Distribution:

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i.e., the k^{th} element of row j = the sum of j consecutive elements on the bottom row of the difference triangle starting at d_{jk}^1 .

For example:

$$\begin{array}{cccc} a_{11} & a_{12} & a_{13} & a_{14} \\ = & 0 & 1 & 4 & 6 \end{array}$$

has difference triangle

$$\begin{array}{ccc} & & 6 \\ & 4 & 5 \\ 1 & 3 & 2 \end{array}$$

row 1 contains the differences between successive antennas

row 2 contains the difference between two antennas separated by one other antenna

- this difference is the sum of the differences between each of the two antennas and the third (separating) one.

The property of a difference triangle which is of interest to us is that:

the sum of the elements in the top $(\frac{n-1}{2})$ rows of a difference triangle
= the sum of the elements in the bottom $(\frac{n-1}{2})$ rows of that difference triangle.

Let $j \leq \frac{n-1}{2}$.

$$\begin{aligned} \sum_{\ell=1}^{n-j} d_{i\ell}^j &= \text{sum of } j^{\text{th}} \text{ row of difference triangle} \\ &= \sum_{\ell=1}^{n-j} \sum_{k=\ell}^{\ell+j-1} d_{ik}^1 \\ &= \sum_{k=1}^{n-1} \alpha_k d_{ik}^1 \end{aligned}$$

$$\begin{aligned} \text{where } \alpha_k &= k \text{ if } 1 \leq k \leq j \\ &= j \text{ if } j \leq k \leq n-j \\ &= n-k \text{ if } n-j \leq k \leq n-1 \end{aligned}$$

$$\begin{aligned} \sum_{\ell=1}^j d_{i\ell}^{n-j} &= \text{sum of } (n-j)^{\text{th}} \text{ row of difference triangle} \\ &= \sum_{\ell=1}^j \sum_{k=\ell}^{\ell+n-j-1} d_{ik}^1 \end{aligned}$$

$$= \sum_{k=1}^{n-1} \beta_k d_{ik}^1$$

$$\text{where } \beta_k = \begin{cases} k & \text{if } 1 \leq k \leq j \\ j & \text{if } j \leq k \leq n-j \\ n-k & \text{if } n-j \leq k \leq n-1 \end{cases}$$

$$\text{Thus } \sum_{\ell=1}^{n-j} d_{i\ell}^j = \sum_{\ell=1}^j d_{i\ell}^{n-j}$$

i.e., sum of j^{th} row = sum of $(n-j)^{\text{th}}$ row for $j=1, \dots, \lfloor \frac{n-1}{2} \rfloor$.

$$\text{Hence } \sum_{k=1}^{\lfloor \frac{n-1}{2} \rfloor} \sum_{\ell=1}^{n-k} d_{i\ell}^k = \sum_{k=1}^{\lfloor \frac{n-1}{2} \rfloor} \sum_{\ell=1}^k d_{i\ell}^{n-k}$$

i.e., sum of bottom $\lfloor \frac{n-1}{2} \rfloor$ rows = sum of top $\lfloor \frac{n-1}{2} \rfloor$ rows.

3. INVERSE-ZOOM ARRAYS

$$\text{Put } p = \lfloor \frac{n-1}{2} \rfloor.$$

There are $\frac{1}{2}p(p+1)$ elements in the top p rows and

$$\frac{1}{2}p(2n-1-p) \text{ elements in the bottom } p \text{ rows.}$$

Since $n \approx 2p$ there are approximately three times as many elements in the bottom half of a difference triangle as in the top half. Since, in general, the bottom half contains the smallest and the top half the largest differences (their sums being equal) it is clear that for every large difference we have several smaller differences. It follows that an 'inverse-zoom' array, in which the large differences outnumber the small ones, cannot be realized except perhaps by having redundancies among the small differences. The calculations which follow make precise the above claim when $n \geq 6$. For $n=4$ and 5 the calculations suggest that perhaps an inverse-zoom array may exist provided it is close to a regular full-fill array (the border-line between zoom arrays and inverse zoom arrays).

Consider the following possibility:

We have m configurations a_{i1}, \dots, a_{in} $i=1, \dots, m$; of n antennas, giving distinct differences (or at least very few redundancies) which make up an inverse-zoom array. That is, the differences get progressively closer together as we go from the small differences to the larger ones.

Put $p = \lfloor \frac{n-1}{2} \rfloor$. Let the average distance between successive differences in the first $S = \frac{m}{2} \lfloor \frac{n-1}{2} \rfloor (2n-1-\lfloor \frac{n-1}{2} \rfloor) = \frac{mp}{2}(2n-1-p)$ differences be f_1 . Let

the largest distance between successive differences in the last $T = \frac{m}{2} p(p+1)$ differences be f_2 . We have $f_1 > f_2$ for an inverse-zoom array.

$$\text{Now } \sum_{i=1}^S f_1 i \leq \text{sum of first } S \text{ differences} \leq \sum_{i=1}^m \sum_{k=1}^p \sum_{\ell=1}^{n-k} d_{i\ell}^k$$

= sum of terms in bottom p rows of the m difference triangles.

Also, sum of terms in top p rows of the m difference triangles

$$= \sum_{i=1}^m \sum_{k=1}^p \sum_{\ell=1}^k d_{i\ell}^{n-k} \leq \text{sum of last } T \text{ differences}$$

$$\leq \sum_{i=1}^T S f_1 + R + f_2 i$$

where $R =$ difference between the $(\frac{n}{2})-T$ difference and the S difference.

$$\text{Thus } \sum_{i=1}^S f_1 i \leq \sum_{i=1}^T S f_1 + R + f_2 i$$

$$\text{Now } \sum_{i=1}^S f_1 i = \frac{1}{2} f_1 S(S+1)$$

$$\text{and } \sum_{i=1}^T S f_1 + R + f_2 i = f_1 S T + R T + \frac{1}{2} f_2 T(T+1) .$$

If n is odd then $n = 2p+1$ and we have

$$S = \frac{m}{2} p(3p+1), \quad T = \frac{m}{2} p(p+1), \quad R = 0 .$$

$$\text{So, } \frac{1}{2} f_1 \frac{m}{2} p(3p+1) (\frac{m}{2} p(3p+1) + 1)$$

$$\leq f_1 \frac{m}{2} p(3p+1) \frac{m}{2} p(p+1) + \frac{1}{2} f_2 \frac{m}{2} p(p+1) (\frac{m}{2} p(p+1) + 1)$$

$$\therefore f_1 \left[\frac{m^2}{8} (9p^4 + 6p^3 + p^2) + \frac{m}{4} (3p^2 + p) \right]$$

$$\leq f_1 \left[\frac{m^2}{4} (3p^4 + 4p^3 + p^2) \right] + f_2 \left[\frac{m^2}{8} (p^4 + 2p^3 + p^2) + \frac{m}{4} (p^2 + p) \right]$$

$$\begin{aligned} \therefore f_1 & \left[\frac{m^2}{8} (3p^4 - 2p^3 - p^2) + \frac{m}{4} (3p^2 + p) \right] \\ & \leq f_2 \left[\frac{m^2}{8} (p^4 + 2p^3 + p^2) + \frac{m}{4} (p^3 + p) \right] . \end{aligned}$$

It follows that for $p \geq 3$ we have $f_2 > f_1$ contradicting our assumption that we have an inverse-zoom array. For $p=2$ and large m we have $f_2 \geq \frac{7}{9} f_1$, which means that for $n=5$ perhaps an inverse-zoom array is possible provided it is not too 'zoomy'.

For $p=1$ ($n=3$) we require $f_2 \geq \frac{f_1}{m}$, so an inverse-zoom array seems possible for large m .

If n is even then $n = 2p+2$ and we have

$$S = \frac{3m}{2} p(p+1), \quad T = \frac{m}{2} p(p+1), \quad R \approx m(p+1) f_2 .$$

$$\begin{aligned} \text{So } \frac{1}{2} f_1 & \frac{3m}{2} p(p+1) \left[\frac{3m}{2} p(p+1) + 1 \right] \\ & \leq f_1 \frac{3m^2}{4} p^2 (p+1)^2 + \frac{m^2}{2} p(p+1)^2 f_2 + \frac{1}{2} f_2 \frac{m}{2} p(p+1) \left[\frac{m}{2} p(p+1) + 1 \right] . \\ \therefore f_1 & \left[\frac{3m^2}{8} (p^4 + 2p^3 + p^2) + \frac{3m}{4} p(p+1) \right] \\ & \leq f_2 \left[\frac{m^2}{8} (p^4 + 6p^3 + 9p^2 + 4p) + \frac{m}{4} p(p+1) \right] . \end{aligned}$$

This contradicts $f_2 < f_1$ if $p \geq 2$.

For $p=1$ ($n=4$) we have (for large m) $f_2 \geq \frac{3}{5} f_1$.

So for $n=4$ an inverse-zoom array may be possible for a small zoom factor.

For $n=2$ anything is possible.