

# Logic gates in Quantum Computing

@CSIRO Co-learning 16/12/21 noon

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\* Intro to QC for astronomers to engineers to management

\* 20+10 min and finance

\* over Zoom

\* at RO

NO TIME FOR CONCEPTS AND IDEAS  
ONLY DETAILS & EQUATIONS

# Logic gates in normal computers

BITS: tiny switches that can be in two configurations  $\{0, 1\}$   
"lightswitch" physical off  $\rightarrow$  0,  $\leftarrow$  on

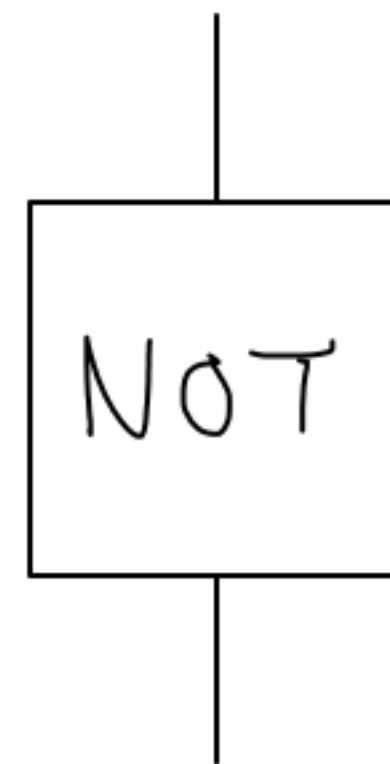
GATES: elementary changes ("operations") applied to bits  
"flip switch"

Example: "NOT" gate  $0 \rightarrow 1$   $1 \rightarrow 0$

Table:

x	not(x)	id(x)	off	on
0	1	0	0	1
1	0	1	0	1

Graphic:



are these all?

NOT MUCH OF INTEREST WITH 1 BIT

# Multi - Bit Gates

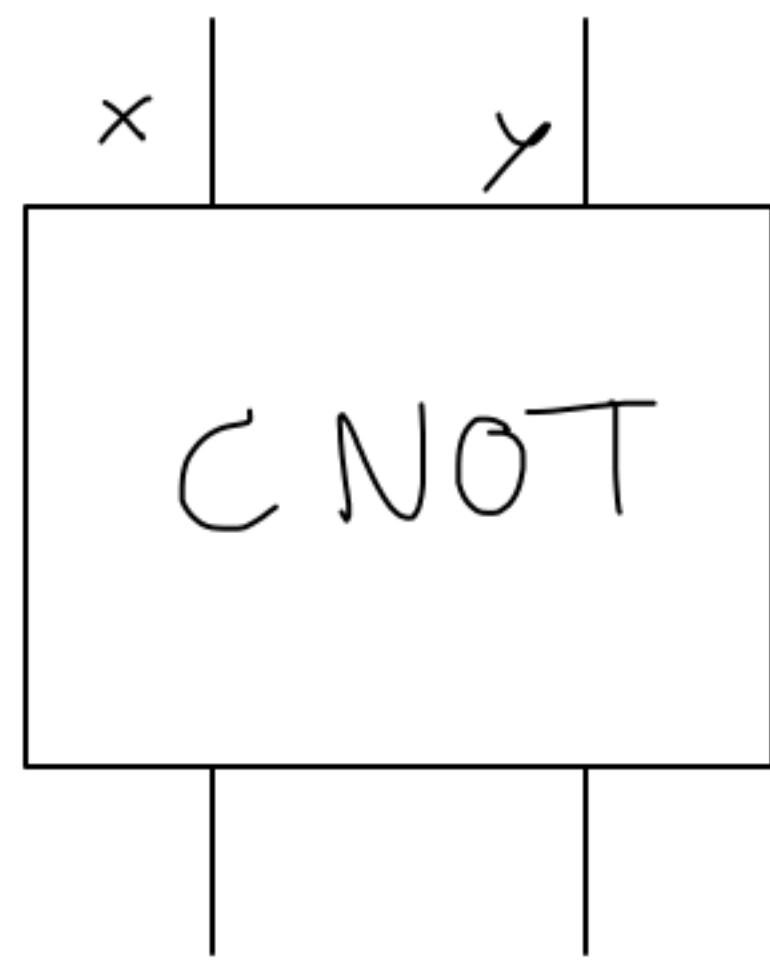
... much more interesting

CNOT :

↪ "control - not"

x	y	CNOT	
0	0	0	0
0	1	0	1
1	0	1	1
1	1	1	0

"if kitchen light on, flip bedroom light"



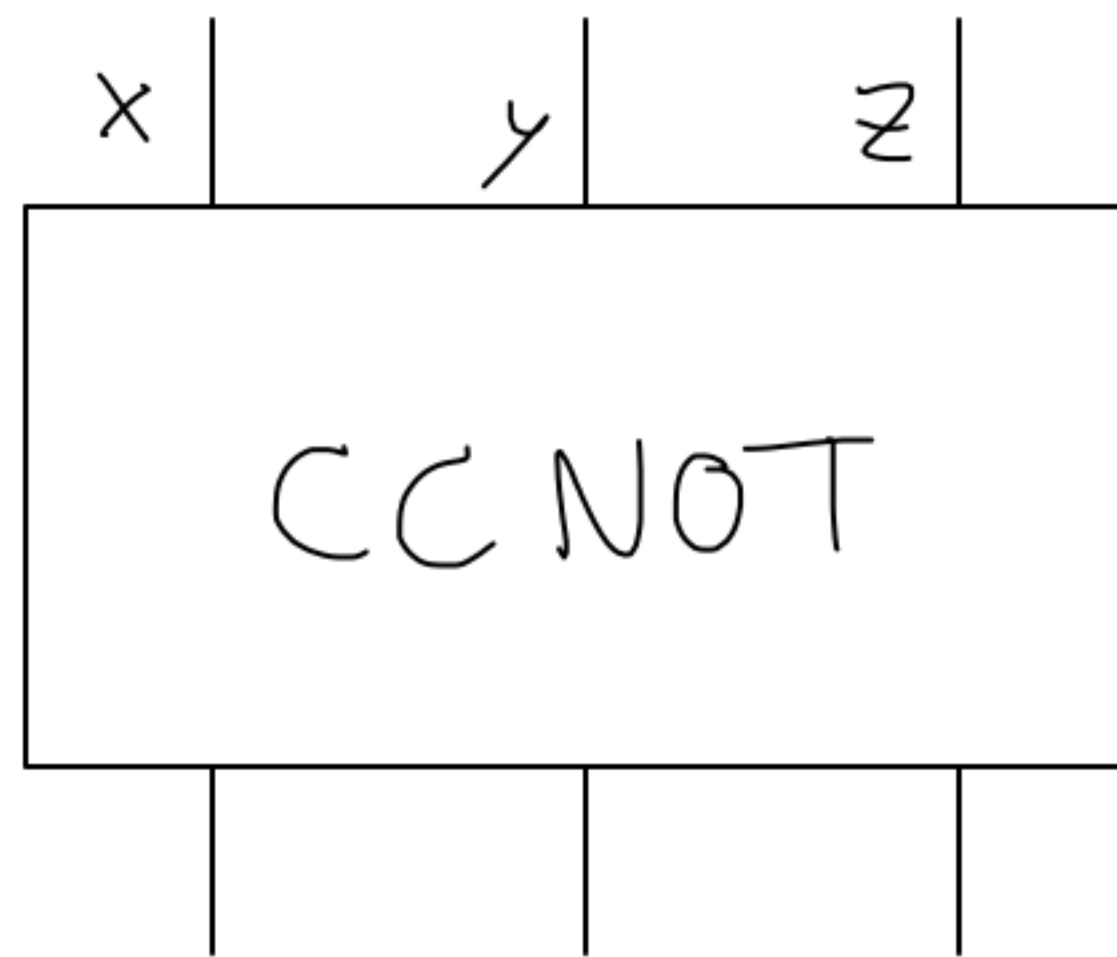
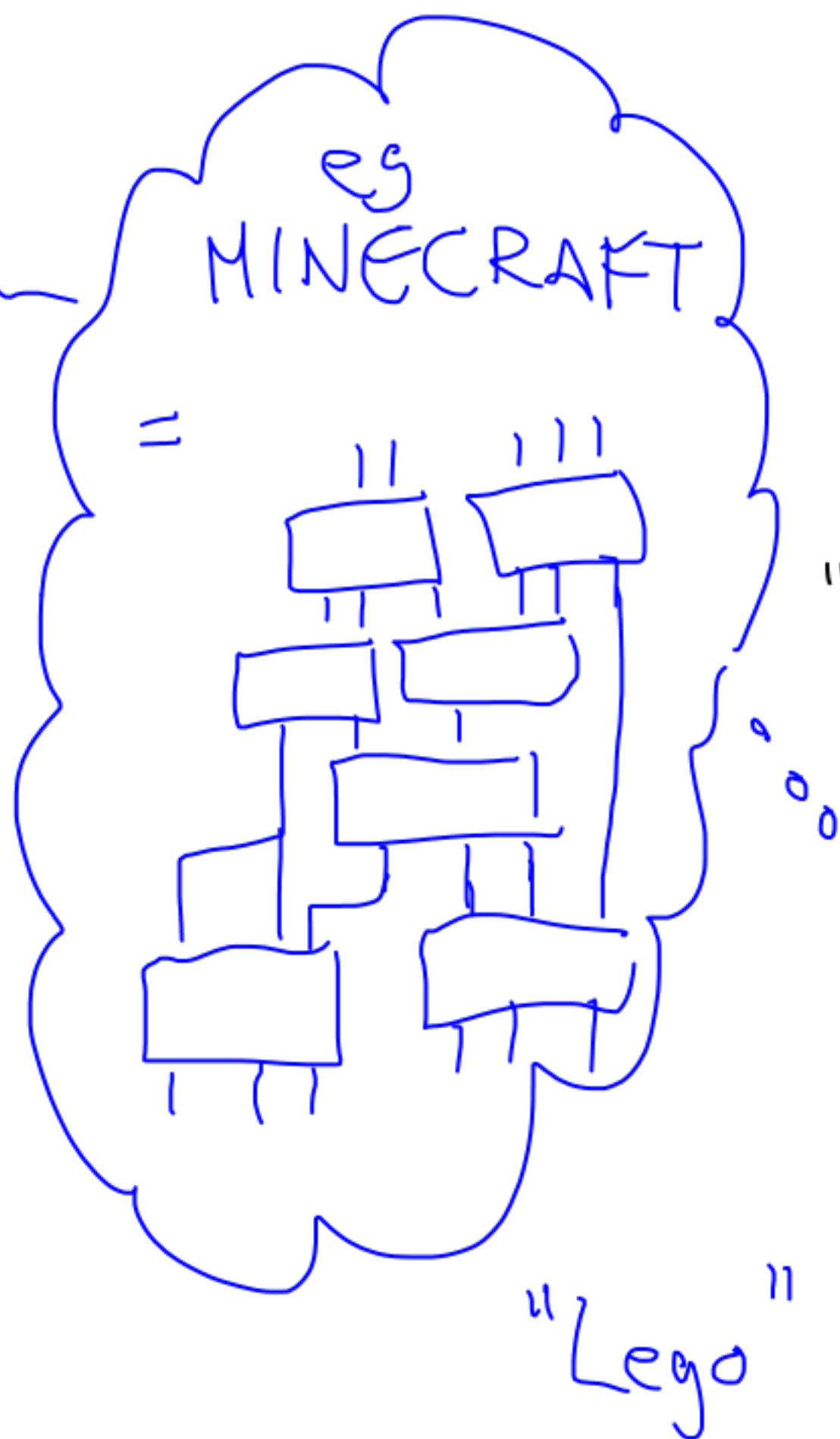
# Multi - Bit Gates

CCNOT :

*controlled - controlled not*

x	y	z	CCNOT
0	0	0	0 0 0
0	0	1	0 0 1
0	1	0	0 1 0
0	1	1	0 1 1
1	0	0	1 0 0
1	0	1	1 0 1
1	1	0	1 1 1
1	1	1	1 1 0

*"if kitchen and bathroom light on, flip bedroom light"*



*"UNIVERSAL":*

ANY SOFTWARE CAN BE CREATED BY CCNOTS

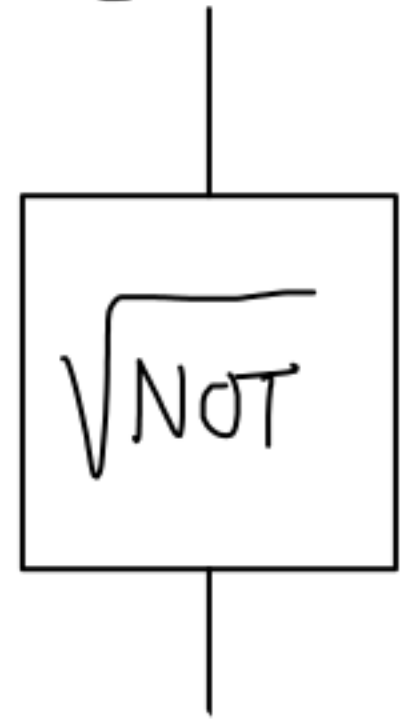
# Logic Quantum Gates

All gates are ultimately physics. Applying gates takes time. What do you get when you stop "half way in" during <sup>"sparks!"</sup> a NOT gate?

A bit is "small, isolated, cold":

**BUT WHAT IS QUANTUM COMPUTING?**

IMPOSSIBLE WITH NORMAL COMPUTERS!



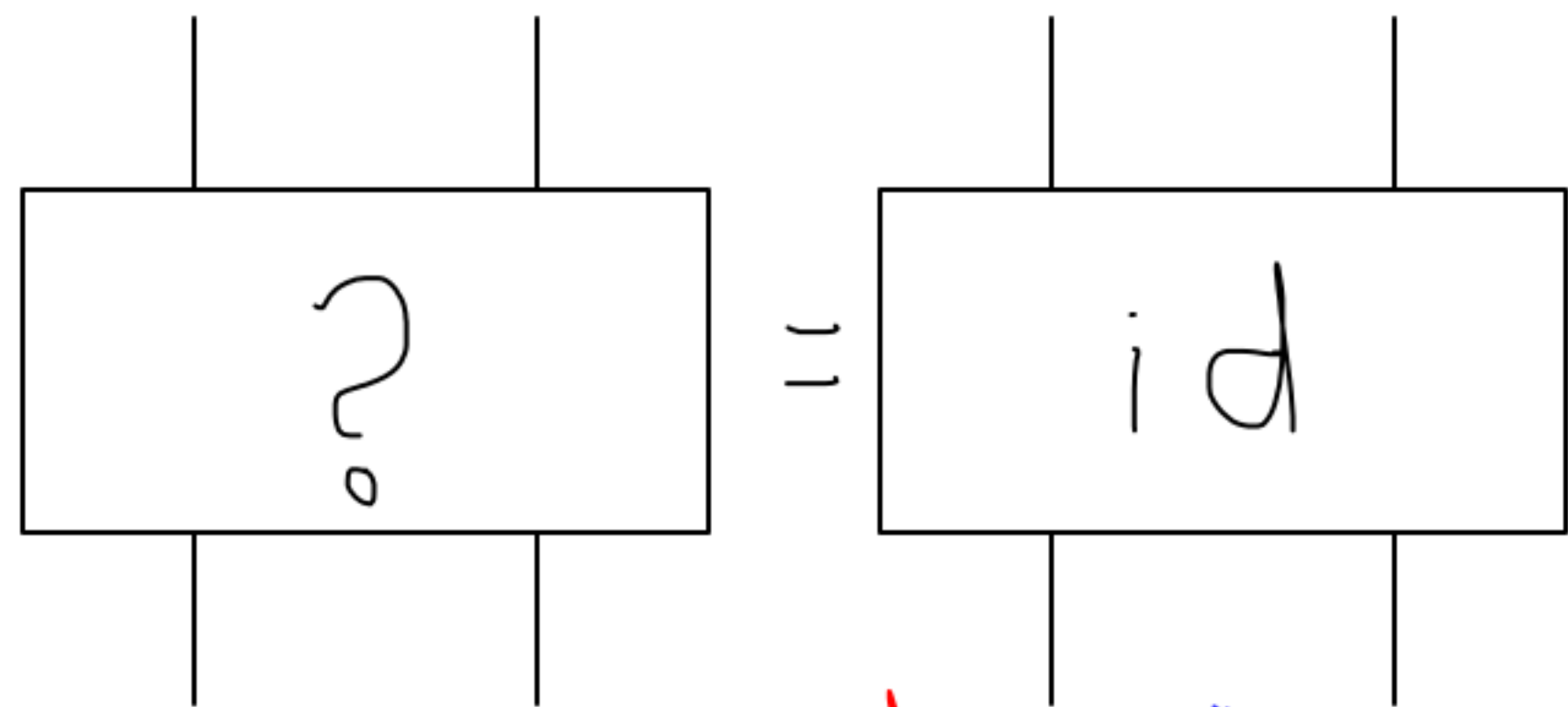
a gate which gives you NOT when applied twice. see at MQ!

**CCNOT +  $\sqrt{\text{NOT}}$  UNIVERSAL FOR QUANTUM COMPUTING**

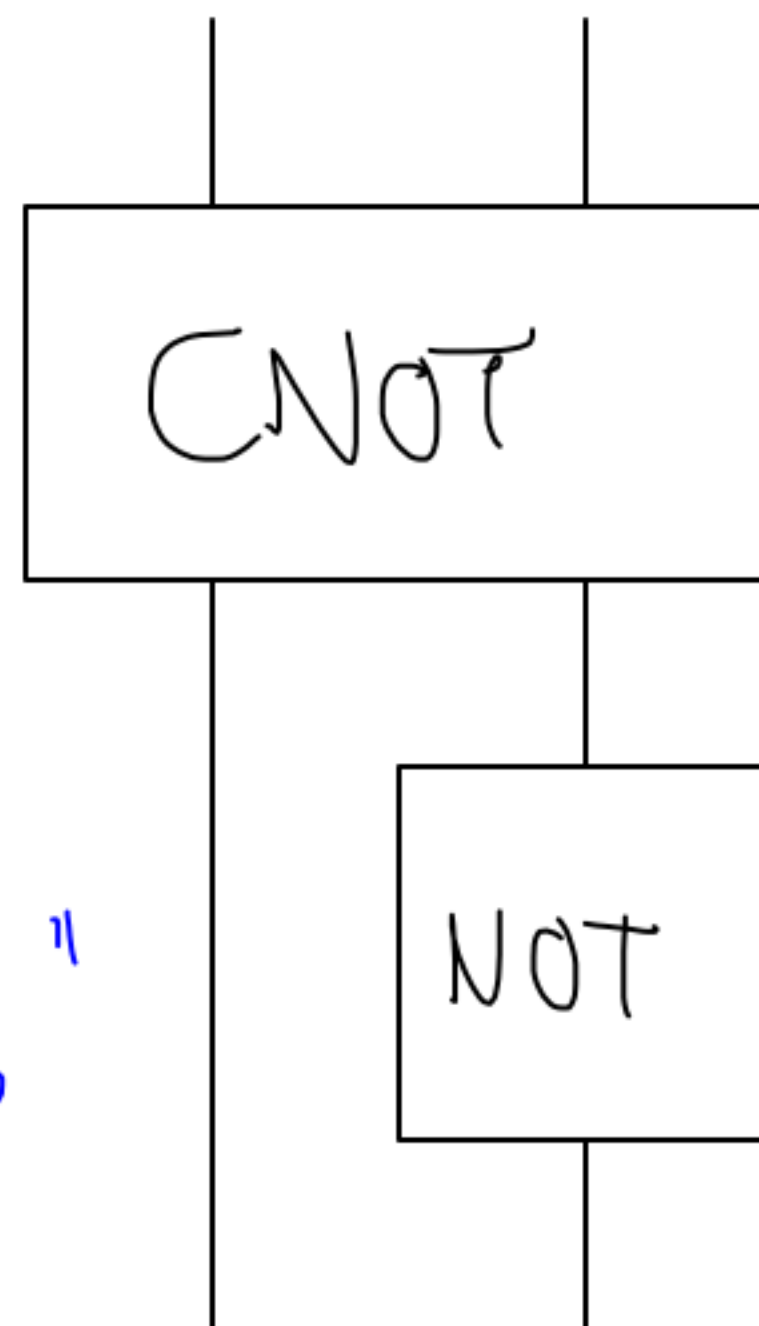
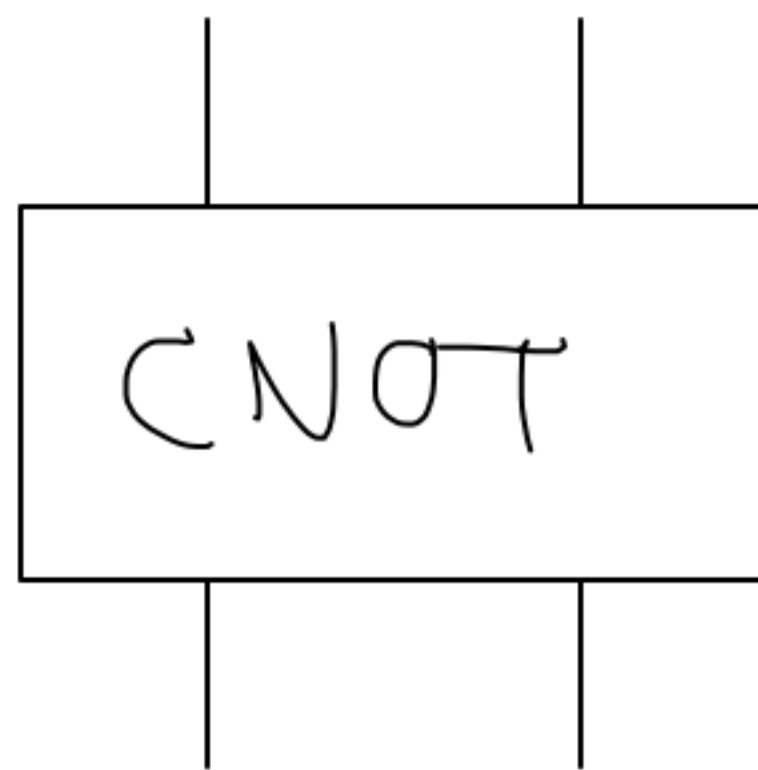
$\text{NOT}^2 = \text{ID}^2 = \text{ID}$   
 $\text{OFF}^2 = \text{OFF}$   
 $\text{ON}^2 = \text{ON}$

A BIT ON WHICH WE CAN APPLY  $\sqrt{\text{NOT}}$  IS CALLED QBIT

mystery box The simplest Quantum Algorithm



IS IT id? "broken"  
FIND OUT FOR SURE



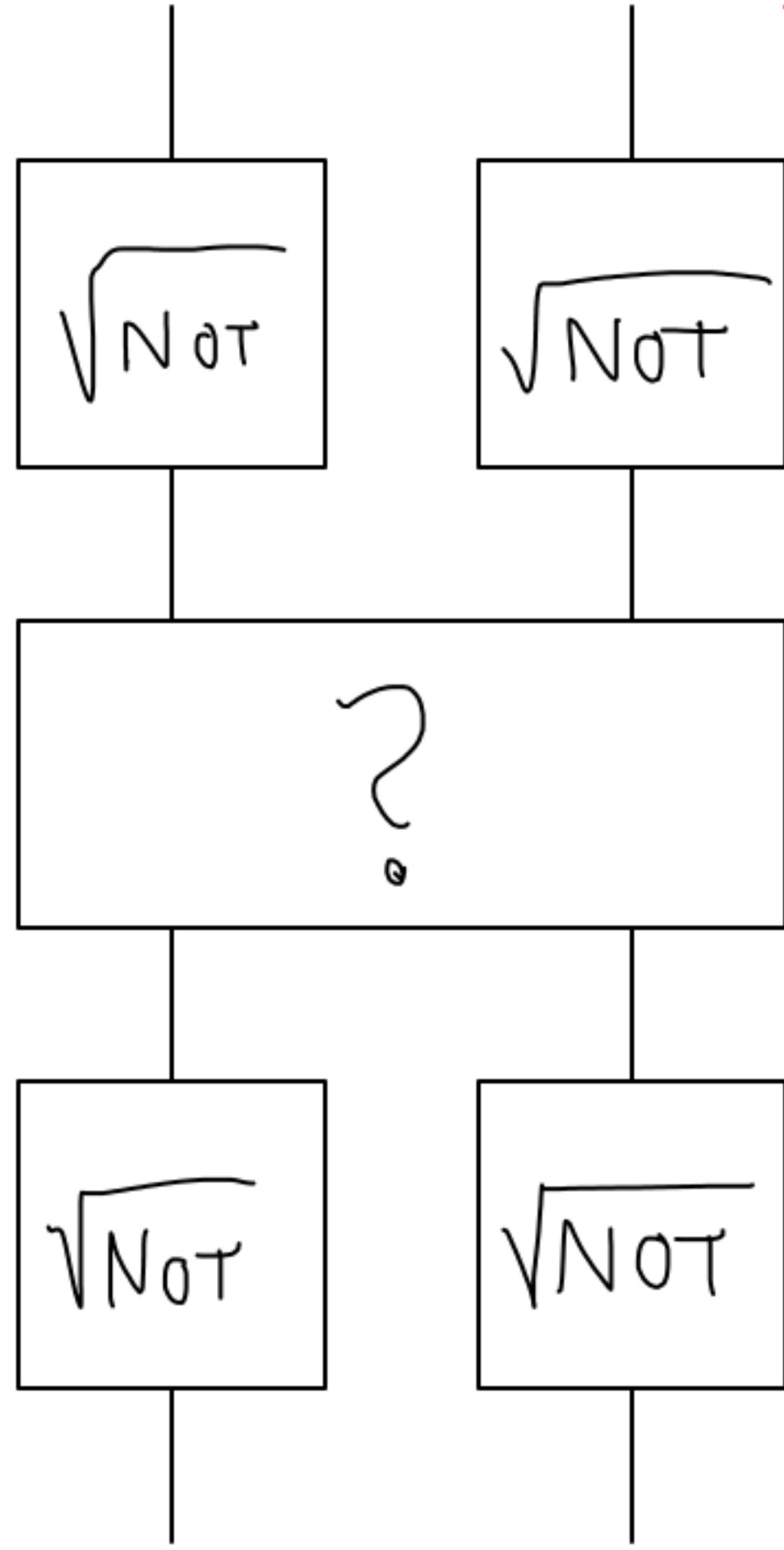
WITH ONE RUN "warranty"

EACH INPUT HAS AMBIGUITIES

x	y	id	CNOT	CNOT NOT
0	0	0 0	0 0	0 1
0	1	0 1	0 1	0 0
1	0	1 0	1 1	1 0
1	1	1 1	1 0	1 1

IMPOSSIBLE WITH  
NORMAL COMPUTER

# A quantum solution



INPUT : 0 1

? = id : 1 0

"broken"

? = CNOT : 0 0

"not broken"

? = CNOT<sub>NOT</sub> : 0 0

⇒ solved !

Quantum algorithms are faster at distinguishing certain things

THIS IDEA GENERALIZES TO EG SHOR

ONLY THE DETAILS...

How does it work?

(Heisenberg)

Idea: describe gates as arrays and bits as lists

$$0 \Rightarrow \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

$$1 \Rightarrow \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$\text{NOT} \Rightarrow \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

array multiplication

$$\text{NOT}(0) = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} = 1$$

EQUIVALENT! BUT MORE MATH FREEDOM:  
CAN TAKE  $\sqrt{\quad}$  OF ARRAY ("experiment")

NOW COMPUTE  
ARRAY SQUARE  
ROOT AND  
MULTIPLY OUT

2 bits: CNOT  $\Rightarrow$

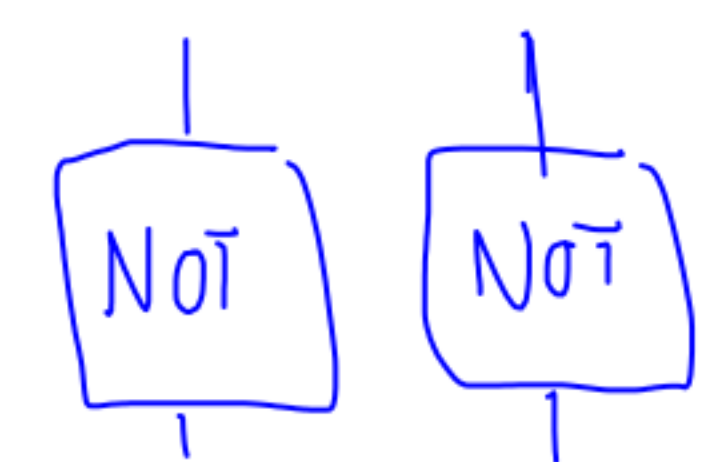
$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

$$00 \Rightarrow \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix}$$

$$01 \Rightarrow \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix}$$

$$10 \Rightarrow \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix}$$

$$11 \Rightarrow \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \end{pmatrix}$$



$$\Rightarrow \begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix}$$



## Conclusions

- \* A quantum computer is a normal one +  $\sqrt{\text{NOT}}$
- \* It is better at distinguishing stuff, sometimes
- \* Understanding it just requires array multiplication

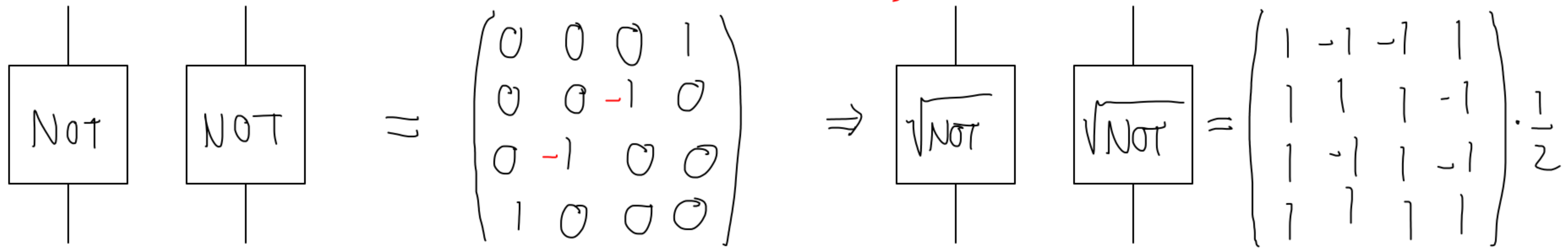
THANKS

PS I cheated :

# Cheating

It turns out that signs are important. We interpret eg.  $\begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix}$  and  $\begin{pmatrix} -1 \\ 0 \\ 0 \\ 0 \end{pmatrix}$  as same.

Gate matrices can also have signs?  
(chosen for convenience; have same effect)



More details

$$\frac{1}{2} \begin{pmatrix} 1 & -1 & -1 & 1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & 1 & 1 \end{pmatrix} \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 0 & 0 & 4 & 0 \\ 0 & -4 & 0 & 0 \\ 4 & 0 & 0 & 0 \end{pmatrix} = \frac{1}{4}$$

